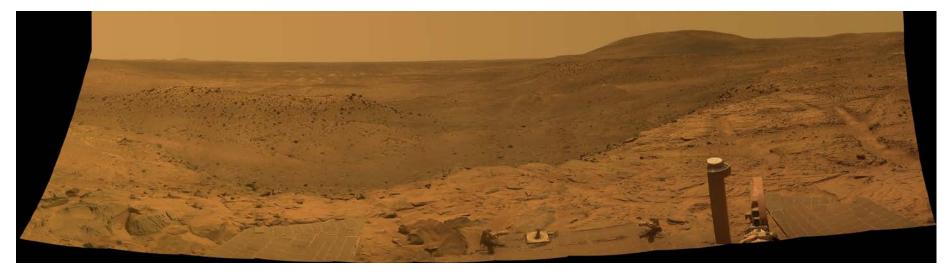


The Design of C++11

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http://www.research.att.com/~bs





Overview

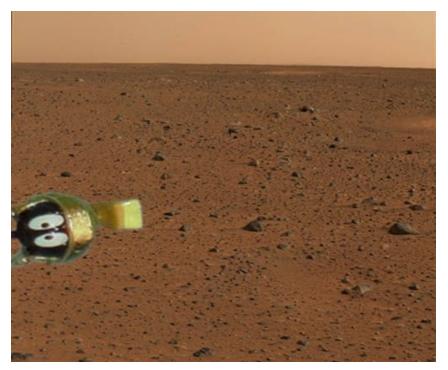
- Aims, Ideals, and history
- C++
- Design rules for C++11
 - With examples
- Case study
 - Concurrency





Programming languages

- A programming language exists to help people express ideas
- Programming language features exist to serve design and programming techniques
- The primary value of a programming language is in the applications written in it

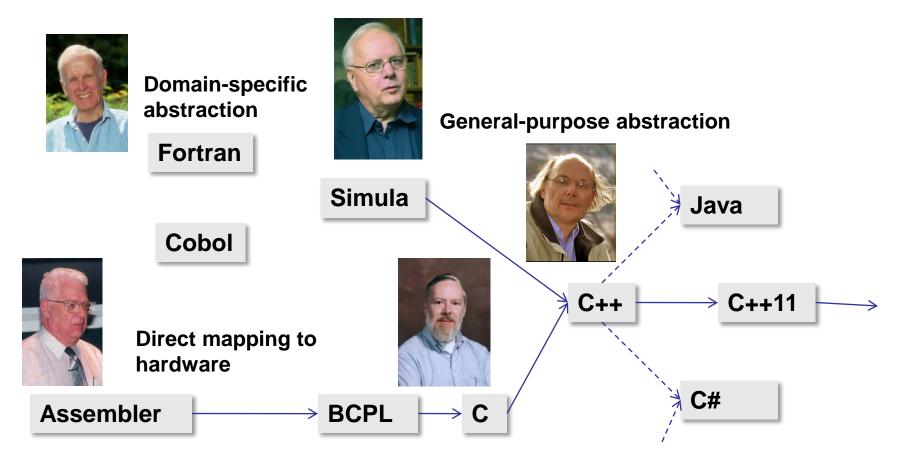


• The quest for better languages has been long and must continue

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Programming Languages



Ideals



- Work at the highest feasible level of abstraction
 - More general, correct, comprehensible, and maintainable code
- Represent
 - concepts directly in code (types, algorithms)
 - independent concepts independently in code
- Represent relationships among concepts directly
 - For example
 - Hierarchical relationships (object-oriented programming)
 - Parametric relationships (generic programming)
- Combine concepts
 - freely
 - but only when needed and it makes sense



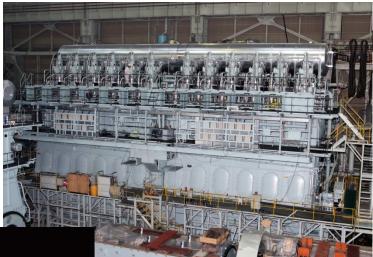
C with Classes –1980

- General abstraction mechanisms to cope with complexity
 From Simula
- General close-to-hardware machine model for efficiency
 - From C
 - Became C++ in 1984
 - Commercial release 1985
 - Non-commercial source license: \$75
 - C++98: ISO standard 1998
 - C++11: 2nd ISO standard 2011





C++ applications









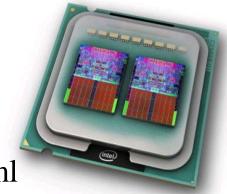






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C++ Applications



www.research.att.com/~bs/applications.html





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C++ Applications









www.lextrait.com/vincent/implementations.html



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C++ ISO Standardization

- Slow, bureaucratic, democratic, formal process
 - "the worst way, except for all the rest"
 - (apologies to W. Churchill)
- About 22 nations (5 to 12 at a meeting)
- Membership have varied
 - 100 to 200+
 - 200+ members currently
 - 40 to 100 at a meeting
 - 70+ currently
- Most members work in industry
- Most members are volunteers
 - Even many of the company representatives
- Most major platform, compiler, and library vendors are represented
 - E.g., IBM, Intel, Microsoft, Sun
- End users are underrepresented





Design?

- Can a committee design?
 - No (at least not much)
 - Few people consider or care for the whole language
- Is C++11 designed
 - Yes
 - Well, mostly: You can see traces of different personalities in C++11
- Committees
 - Discuss
 - Bring up problems
 - "Polish"
 - Are brakes on innovation

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Overall goals for C++11

- Make C++ a better language for systems programming and library building
 - Rather than providing specialized facilities for a particular subcommunity (e.g. numeric computation or Windows-style application development)
 - Build directly on C++'s contributions to systems programming



- Make C++ easier to teach and learn
 - Through increased uniformity, stronger guarantees, and facilities supportive of novices (there will always be more novices than experts)

C++11



- C++11 is not science fiction
 - Became an ISO Standard in 2011
 - Every feature is implemented somewhere
 - And shipping, e.g. Microsoft, GCC, Clang, EDG, ...
 - E.g. GCC 4.7: Rvalues, Variadic templates, Initializer lists, Static assertions, **auto**, New function declarator syntax, Lambdas, Right angle brackets, Extern templates, Strongly-typed **enum**s, **constexpr**, Delegating constructors (patch), Raw string literals, Defaulted and deleted functions, **noexcept**, Local and unnamed types as template arguments, range-**for**, user-defined literals, ...
 - Standard library components are shipping widely
 - E.g. GCC, Microsoft, Boost



Rules of thumb / Ideals

- Integrating features to work in combination is the key
 - And the most work
 - The whole is much more than the simple sum of its part
- Maintain stability and compatibility
- Prefer libraries to language extensions
- Prefer generality to specialization
- Support both experts and novices
- Increase type safety
- Improve performance and ability to work directly with hardware
- Make only changes that change the way people think
- Fit into the real world

Maintain stability and compatibility

- "Don't break my code!"
 - There are billions of lines of code "out there"
 - There are millions of C++ programmers "out there"
- "Absolutely no incompatibilities" leads to ugliness
 - We introduce new keywords as needed: auto (recycled), decltype, constexpr, thread_local, nullptr
 - Example of incompatibility: static_assert(4<=sizeof(int),"error: small ints");</p>





Support both experts and novices

- *Example*: minor syntax cleanup
 vector<list<int>> v; // note the "missing space"
- *Example*: simplified iteration for (auto x : v) cout << x <<'\n';
- *Note*: Experts don't easily appreciate the needs of novices
 - Example of what we couldn't get just now

string s = "12.3";

double x = lexical_cast<double>(s); // extract value from string



Uniform initialization

 You can use {}-initialization for all types in all contexts int a[] = { 1,2,3 }; vector<int> v { 1,2,3};

```
vector<string> geek_heros = {
    "Dahl", "Kernighan", "McIlroy", "Nygaard ", "Ritchie", "Stepanov"
};
```

complex<double> z{1,2}; // invokes constructor
struct S { double x, y; } s {1,2}; // no constructor (just initialize members)



Uniform initialization

- {}-initialization X{v} yields the same value of X in every context X x{a};
 - X* p = new X{a};
 - $\mathbf{z} = \mathbf{X}\{\mathbf{a}\};$ // use as cast

```
void f(X);
```

f({a}); *// function argument (of type X)*

```
X g() {
```

// ...

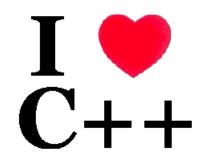
return {a}; // function return value (function returning X)
}

Y::Y(a): X{a} { /* ... */ } // base class initializer Stroustrup - Wroclaw'12



Uniform initialization

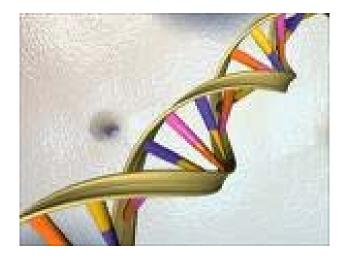
 {}-initialization does not narrow int x1 = 7.9; // x1 becomes 7 int x2 {7.9}; // error: narrowing conversion





Prefer libraries to language extensions

- Libraries deliver more functionality
- Libraries are immediately useful
- *Problem*: Enthusiasts prefer language features
 - see library as 2nd best
- *Example*: New library components
 - std::thread, std::future, ...
 - Threads ABI; not thread built-in type
 - std::unordered_map, std::regex, ...
 - Not built-in associative array





Prefer generality to specialization

- *Example*: Prefer improvements to abstraction mechanisms over separate new features
 - Inherited constructor

```
template<class T> class Vector : std::vector<T> {
    using vector::vector<T>; // inherit all constructors
    // ...
    };
    Move semantics supported by rvalue references
    template<class T> class vector {
        // ...
        void push_back(T&& x); // move x into vector
        // avoid copy if possible
```

- };
- *Problem*: people love small isolated features

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Not a reference

ł



Move semantics

 Often we don't want two copies, we just want to move a value vector<int> make_test_sequence(int n)

```
vector<int> res;
for (int i=0; i<n; ++i) res.push_back(rand_int());
return res; // move, not copy
}
```

```
vector<int> seq = make_test_sequence(1000000); // no copies
```

- New idiom for arithmetic operations:
 - Matrix operator+(const Matrix&, const Matrix&);
 - a = b+c+d+e; // no copies

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Move semantics

• Move constructor template<typename T> class vector { // ... vector(vector&& v) { **elem = v.elem;** // "steal" v's representation SZ = V.SZ;elem = nullptr; // leave v empty sz = 0; private: T* elem; int sz; Stroustrup - Wroclaw'12



Increase type safety

- Approximate the unachievable ideal
 - Example: Strongly-typed enumerations enum class Color { red, blue, green };
 - int x = Color::red; // error: no Color->int conversion
 - Color y = 7;
 - Color z = red;

- *Il error: no color->thi conversion Il error: no int->Color conversion Il error: red not in scope Il fine*
- *Example*: Support for general resource management
 - **std::unique_ptr** (for ownership)
 - **std::shared_ptr** (for sharing)
 - Garbage collection ABI

Color c = Color::red;

Improve performance and the ability to work of the solution of

- Embedded systems programming is very important
 - *Example*: address array/pointer problems
 - array<int,7> s; // fixed-sized array
 - *Example*: Generalized constant expressions (think ROM)
 constexpr int abs(int i) { return (0<=i) ? i : -i; } // can be constant expression</p>

struct Point { // "literal type" can be used in constant expression
int x, y;

constexpr Point(int xx, int yy) : x{xx}, y{yy} { }
};

constexpr Point p1{1,2}; // must be evaluated at compile time: ok constexpr Point p2{1,abs(x)}; // ok?: is x is a constant expression? Stroustrup - Wroclaw'12



Make only changes that change the way people think

- Think/remember:
 - Object-oriented programming
 - Generic programming
 - Concurrency
 - ...
- But, most people prefer to fiddle with details
 - So there are dozens of small improvements
 - All useful somewhere
 - long long, static_assert, raw literals, thread_local, unicode types, ...
 - *Example*: A null pointer keyword
 - void f(int);
 - void f(char*);
 - **f(0);**

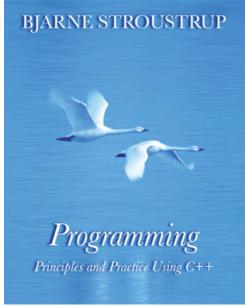
// call f(int); f(nullptr); // call f(char*);

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Fit into the real world

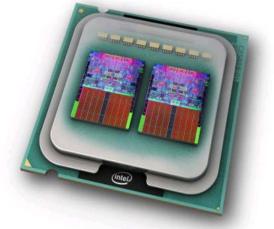
- *Example*: Existing compilers and tools must evolve
 - Simple complete replacement is impossible
 - Tool chains are huge and expensive
 - There are more tools than you can imagine
 - C++ exists on many platforms
 - So the tool chain problems occur N times
 (for each of M tools)
- *Example*: Education
 - Teachers, courses, and textbooks
 - Often mired in 1970s thinking ("C is the perfect language")
 - Often mired in 1980s thinking ("OOP: Rah! Rah!! Rah!!!")
 - "We" haven't completely caught up with C++98!
 - "legacy code breeds more legacy code"





Areas of language change

- Machine model and concurrency Model
 - Threads library (std::thread)
 - Atomics ABI
 - Thread-local storage (thread_local)
 - Asynchronous message buffer (std::future)
- Support for generic programming
 - (no concepts $\ensuremath{\mathfrak{S}}$)
 - uniform initialization
 - auto, decltype, lambdas, template aliases, move semantics, variadic templates, range-for, ...
- Etc.
 - static_assert
 - improved **enum**s
 - long long, C99 character types, etc.





Standard Library Improvements

- New containers
 - Hash Tables (**unordered_map**, etc.)
 - Singly-linked list (forward_list)
 - Fixed-sized array (array)
- Container improvements
 - Move semantics (e.g. push_back)
 - Intializer-list constructors
 - Emplace operations
 - Scoped allocators
- More algorithms (just a few)
- Concurrency support
 - thread, mutex, lock, ...
 - future, async, ...
 - Atomic types
- Garbage collection ABI





Standard Library Improvements

- Regular Expressions (**regex**)
- General-purpose Smart Pointers (unique_ptr, shared_ptr, ...)
- Extensible Random Number Facility
- Enhanced Binder and function wrapper (bind and function)
- Mathematical Special Functions
- Tuple Types (**tuple**)
- Type Traits (lots)





Template What is C++? meta-programming!

A hybrid language

Buffer overflows

Too big!



A multi-paradigm programming language

It's C!

Embedded systems programming language

Supports generic programming

Low level!

An object-oriented programming language

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A random collection of features 32

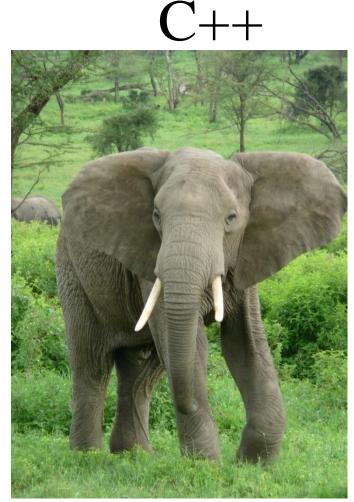
C++11



- It *feels* like a new language
 - Compared to C++98
- It's not just "object oriented"
 - Many of the key user-defined abstractions are not objects
 - Types
 - Classifications and manipulation of types (types of types)
 - I miss "concepts"
 - Algorithms (generalized versions of computation)
 - Resources and resource lifetimes
- The pieces fit together much better than they used to







A light-weight abstraction programming language

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So, what does "light-weight abstraction" mean?"

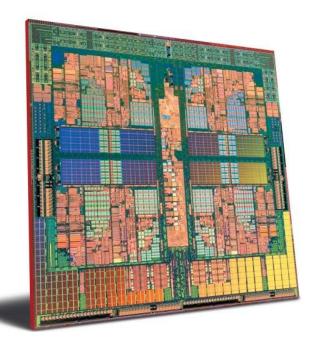
- The design of programs focused on the design, implementation, and use of abstractions
 - Often abstractions are organized into libraries
 - So this style of development has been called "library-oriented"
- C++ emphasis
 - Flexible static type system
 - Small abstractions
 - Performance (in time and space)
 - Ability to work close to the hardware





Case study

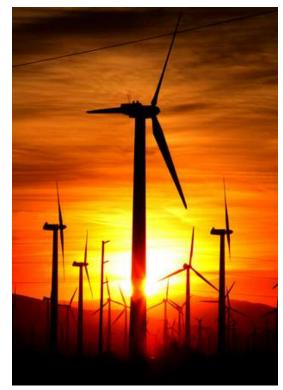
- Concurrency
 - "driven by necessity"
 - More than ten years of experience





Case study: Concurrency

- What we want
 - Ease of programming
 - Writing correct concurrent code is hard
 - Portability
 - Uncompromising performance
 - System level interoperability
- We can't get everything
 - No one concurrency model is best for everything
 - De facto: we can't get all that much
 - "C++ is a systems programming language"
 - (among other things) implies serious constraints





Concurrency overview

- Foundation
 - Memory model
 - atomics
- Concurrency library components
 - std::thread
 - std::mutex (several)
 - std::lock (several)
 - std::condition (several)
 - std::future, std::promise, std::packaged_task
 - std::async()
- Resource management
 - std::unique_ptr, std::shared_ptr
 - GC ABI Stroustrup Wroclaw'12





Memory model

• A memory model is an agreement between the machine architects and the compiler writers to ensure that most programmers do not have to think about the details of modern computer hardware.

ll thread 1:	<i>II thread 2:</i>
char c;	char b;
c = 1;	b = 1;
int x = c;	int $y = b$;

x==1 and y==1 as anyone would expect
(but don't try that for two bitfields of the same word)



Memory model

- Two threads of execution can update and access separate memory locations without interfering with each other.
- But what is a "memory location?"
 - A memory location is either an object of scalar type or a maximal sequence of adjacent bit-fields all having non-zero width.
 - For example, here S has exactly four separate memory locations:
 struct S {

char a;	<i> location #1</i>
int b:5;	<i>II location #2</i>
unsigned c:11;	
unsigned :0;	<pre>// note: :0 is "special"</pre>
unsigned d:8;	<i>II location #3</i>
<pre>struct {int ee:8;} e;</pre>	<i>II location #4</i>
};	Stroustrup - Wroclaw'12



Atomics ("here be dragons!")

- Components for fine-grained atomic access
 - provided via operations on atomic objects (in <cstdatomic>)
 - Low-level, messy, and shared with C (making the notation messy)
 - what you need for lock-free programming
 - what you need to implement std::thread, std::mutex, etc.
 - Several synchronization models, CAS, fences, ...

```
enum memory_order { // regular (non-atomic) memory synchronization order
memory_order_relaxed, memory_order_consume, memory_order_acquire,
memory_order_release, memory_order_acq_rel, memory_order_seq_cst
```

};

C atomic_load_explicit(const volatile A* object, memory_order);

void atomic_store_explicit(volatile A *object, C desired, memory_order order);

bool atomic_compare_exchange_weak_explicit(volatile A* object, C * expected, C
 desired, memory_order success, memory_order failure);

II ... lots more ...





Threading

- You can
 - wait for a thread for a **specified time**
 - control access to some data by **mutual exclusion**
 - control access to some data using locks
 - wait for an action of another task using a **condition variable**
 - return a value from a thread through a **future**



Concurrency: std::thread

#include<thread>

```
void f() { std::cout << "Hello "; }</pre>
```

```
struct F {
    void operator()() { std::cout << "parallel world "; }
};</pre>
```

```
int main()
```

{

```
std::thread t1{f}; // f() executes in separate thread
std::thread t2{F()}; // F()() executes in separate thread
} // spot the bugs
```

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Concurrency: std::thread

```
int main()
```

{

}

```
std::thread t1{f}; // f() executes in separate thread
std::thread t2{F()}; // F()() executes in separate thread
```

t1.join();	<i> wait for t1</i>
t2.join();	ll wait for t2

// and another bug: don't write to cout without synchronization



Thread – pass result (primitive)

void f(vector<double>&, double* res); // place result in res
struct F {

```
vector& v; double* res;
```

 $F(vector < double > \& vv, double * p) : v\{vv\}, res\{p\} \{ \}$

```
void operator()(); // place result in res
```

```
};
```

```
int main()
```

```
{
```

}

```
double res1; double res2;
```



Thread – pass argument and result

double* f(const vector<double>& v); // read from v return result
double* g(const vector<double>& v); // read from v return result

void user(const vector<double>& some_vec) // note: const

{

}

```
double res1, res2;
thread t1 {[&]{ res1 = f(some_vec); }}; // lambda: leave result in res1
thread t2 {[&]{ res2 = g(some_vec); }}; // lambda: leave result in res2
// ...
t1.join();
t2.join();
cout << res1 << ' ' << res2 << '\n';</pre>
```



No cancellation/interruption

- When a **thread** goes out of scope the program is **terminate**()d unless its task has completed. That's obviously to be avoided.
- There is no way to request a **thread** to terminate (i.e. request that it exit as a soon as possible and as gracefully as possible) or to force a thread to terminate (i.e. kill it). We are left with the options of
- designing our own cooperative ``interruption mechanism" (with a piece of shared data that a caller thread can set for a called thread to check (and quickly and gracefully exit when it is set)),
- ``going native" (using **thread::native_handle**() to gain access to the operating system's notion of a thread),
- kill the process (std::quick_exit()),
- kill the program (std::terminate().



Mutual exclusion: std::mutex

- A **mutex** is a primitive object use for controlling access in a multi-threaded system.
- A **mutex** is a shared object (a resource)
- Simplest use: std::mutex m; int sh; // shared data // ... m.lock(); // manipulate shared data: sh+=1;

m.unlock();





Mutex – try_lock()

• Don't wait unnecessarily

```
std::mutex m;
```

```
int sh; // shared data
// ...
if (m.try_lock()) { // manipulate shared data:
    sh+=1;
    m.unlock();
else {
    // maybe do something else
}
```



Mutex - try_lock_for()

 Don't wait for too long: std::timed_mutex m; int sh; // shared data

// ...

```
if (m.try_lock_for(std::chrono::seconds(10))) { // Note: time
    // manipulate shared data:
    sh+=1;
    m.unlock();
}
else {
    // we didn't get the mutex; do something else
}
```



Mutex - try_lock_until()

• We can wait until a fixed time in the future: std::timed_mutex m;

int sh; // shared data

// ...

if (m.try_lock_until(midnight)) { // manipulate shared data:
 sh+=1;
 m.unlock();

} else {

}

II we didn't get the mutex; do something else



Recursive mutex

• In some important use cases it is hard to avoid recursion std::recursive_mutex m; int sh; // shared data // ... void f(int i) { // ... m.lock(); *II manipulate shared data:* **sh**+=1; if (--i>0) f(i);

```
m.unlock();
```



RAII for mutexes: std::lock

• A lock represents local ownership of a non-local resource (the **mutex**)

mutex m;

int sh; // shared data

```
void f()
{
    // ...
    unique_lock<mutex>lck(m);
    // manipulate shared data:
    sh+=1;
} // implicitly release the mutex
```

Il grab (acquire) the mutex



Potential deadlock

• Unstructured use of multiple locks is hazardous:

mutex m1; mutex m2; int sh1; // shared data int sh2; // ... void f() { // ...



```
unique_lock<mutex> lck1(m1);
unique_lock<mutex> lck2(m2);
// manipulate shared data:
sh1+=sh2;
```



RAII for mutexes: std::lock

We can safely use several locks
 void f() {

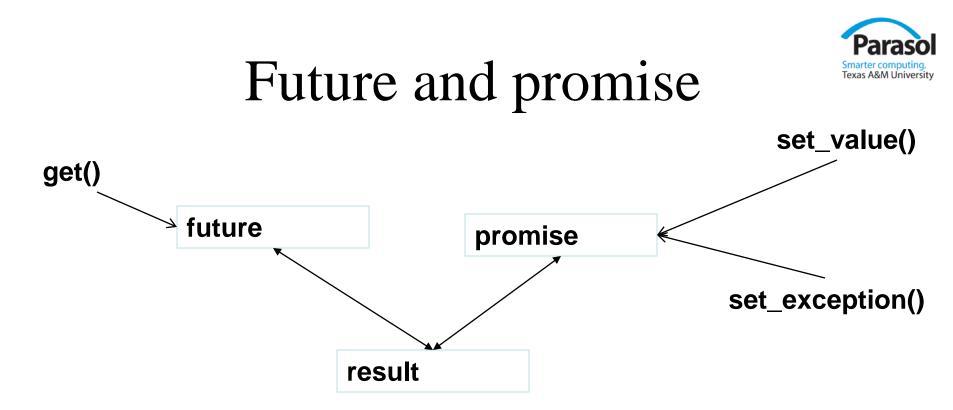
// ...

```
unique_lock<mutex> lck2(m2,defer_lock);
unique_lock<mutex> lck3(m3,defer_lock);
// ...
```

```
lock(lck1,lck2,lck3);
```

// manipulate shared data

} || implicitly release the mutexes



- future+promise provides a simple way of passing a value from one thread to another
 - No explicit synchronization
 - Exceptions can be transmitted between threads



Future and promise

- Get from a future<X> called f:
 X v = f.get();// if necessary wait for the value to get
- Put to a promise<X> called p (attached to f): try {
 - X res;
 - *Il compute a value for res*
 - p.set_value(res);
 - } catch (...) {

}

```
// oops: couldn't compute res
p.set_exception(std::current_exception());
```

async() – pass argument and return result

double* f(const vector<double>& v); // read from v return result
double* g(const vector<double>& v); // read from v return result

```
void user(const vector<double>& some_vec)  // note: const
{
    auto res1 = async(f,some_vec);
    auto res2 = async(g,some_vec);
    // ...
    cout << *res1.get() << ' ' << *res2.get() << '\n'; // futures
}</pre>
```

- Much more elegant than the explicit thread version
 - And most often faster

async()



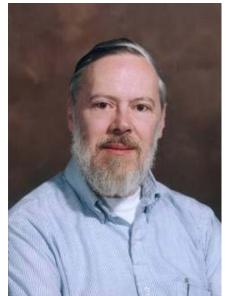
Simple launcher using the variadic template interface
 template<class T, class V> struct Accum { /* accumulator function object };

```
void comp(vector<double>& v) // spawn many
{
    auto b = &v[0];
    auto sz = v.size();
```

```
auto f0 = async(Accum, b, b+sz/4, 0.0);auto f1 = async(Accum, b+sz/4, b+sz/2, 0.0);auto f2 = async(Accum, b+sz/2, b+sz*3/4, 0.0);auto f3 = async(Accum, b+sz*3/4, b+sz, 0.0);return f0.get()+f1.get()+f2.get()+f3.get();
```

}





Thanks!

- C and Simula
 - Brian Kernighan
 - Doug McIlroy
 - Kristen Nygaard
 - Dennis Ritchie
 - ..
- ISO C++ standards committee
 - Steve Clamage
 - Francis Glassborow
 - Andrew Koenig
 - Tom Plum
 - Herb Sutter
 - —
- C++ compiler, tools, and library builders
 - Beman Dawes
 - David Vandevoorde
- Application builders Stroustrup - Wroclaw'12





Parasol Smarter computing. Texas A&M University

More information

- My home pages
 - C++11 FAQ
 - Papers, FAQs, libraries, applications, compilers, ...
 - Search for "Bjarne" or "Stroustrup"
 - "Software Development for Infrastructure" paper
 - My HOPL-II and HOPL-III papers
- The Design and Evolution of C++ (Addison Wesley 1994)
- The ISO C++ standard committee's site:
 - All documents from 1994 onwards
 - Search for "WG21"
- The Computer History Museum
 - Software preservation project's C++ pages
 - Early compilers and documentation, etc.
 - http://www.softwarepreservation.org/projects/c_plus_plus/
 - Search for "C++ Historical Sources Archive"

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