## Course $C^{++}$, Exercise List 7

Deadline: 28.04.2016

In this exercise list, we implement matching, using the trees of exercise list 5, and std: :vector.

1. Write bool operator $==$ ( tree t 1 , tree t 2 ), that returns true iff the trees t 1 and t 2 are equal. One can write a recursive procedure, but we won't do that, because we want to use std: :vector.
We will use the following definition:

$$
\begin{array}{ll}
S \cup\left\{f\left(t_{1}, \ldots, t_{n}\right)\right\} \equiv\left\{g\left(u_{1}, \ldots, u_{m}\right)\right\} & \Rightarrow \quad \text { false if } n \neq m, \text { or } f \neq g \\
S \cup\left\{f\left(t_{1}, \ldots, t_{n}\right)\right\} \equiv\left\{f\left(u_{1}, \ldots, u_{n}\right)\right\} & \Rightarrow \quad S \cup\left\{t_{1} \equiv u_{1}, \ldots, t_{n} \equiv u_{n}\right\} \\
\emptyset & \Rightarrow \quad \text { true. }
\end{array}
$$

If one wants to know whether trees $t_{1}, t_{2}$ are equal, one starts with $S=$ $\left\{t_{1} \equiv t_{2}\right\}$.
Use std::vector< std::pair< tree, tree>> to implement the set of equations. Take the pair at the end, process it, and if necessary put new pairs back at the end.
2. Download matching.h from the course homepage.

Complete operator ( ) (const tree\& t) const in struct matching. It is very similar to the substitution function of List 5 .
Use of replacesubtree, in order to avoid unecessary copying. Make sure that the non-const versions of functor ( ) and operator [] ( ) are removed in class tree.
3. A tree $t_{1}$ can be matched into a tree $t_{2}$, if there exists a substitution $\Theta$, s.t. $t_{1} \Theta=t_{2}$.

Matching can be implemented recursively, but we interested in using std: :list, or std: :vector, so we (=you) will implement matching using the following definition:

$$
\begin{array}{ll}
\left(\Theta, S \cup\left\{f\left(t_{1}, \ldots, t_{n}\right) / g\left(u_{1}, \ldots, u_{m}\right)\right\}\right) & \Rightarrow \text { false if } n \neq m, \text { or } f \neq g \\
\left(\Theta, S \cup\left\{f\left(t_{1}, \ldots, t_{n}\right) / f\left(u_{1}, \ldots, u_{n}\right)\right\}\right) & \Rightarrow\left(\Theta, S \cup\left\{t_{1} / u_{1}, \ldots, t_{n} / u_{n}\right\}\right) \\
(\Theta, S \cup\{V / t\}) & \Rightarrow(\Theta \cup\{V:=t\}, S) \text { if } V \Theta \text { is undefined, } \\
(\Theta, S \cup\{V / t\}) & \Rightarrow \text { false if } V \Theta \text { is defined, and } V \Theta \neq t \\
(\Theta, S \cup\{V / t\}) & \Rightarrow(\Theta, S) \text { if } V \Theta \text { is defined, and } V \Theta=t .
\end{array}
$$

We assume that $f, g$ are functors that are not variables. $V$ is a tree without subterms with a functor that is a variable. In order to decide whether a functor is a variable, use matching: :isvariable( const std::string\& s ). (A functor is a variable if it starts with an underscore.)
$t$ is an arbitrary tree.
States have form $(\Theta, S)$, where $\Theta$ is the current matching and $S$ contains the pairs of trees to be matched. Matching starts with ( $\emptyset,\{$ from/into $\}$ ). Use a matching and std::vector< std::pair<tree,tree>>.
There is a complication that matching may fail. Since this is not exceptional, one should not throw an exception in this case. We solve the problem by returning a std: :list<matching>, which is empty, when matching fails, and contains one matching, when matching succeeds.

Below are some examples:

```
f( _X, _Y ) into f(f(a), f(b) )
    _X := f(a) _Y := f(b)
f( _X, _X ) into f(f(a), f(b) )
        fails
f( _X, _Y ) into g( f(a), f(b) )
        fails
f( _X, _X ) into f( f(a), f(a) )
    X := f(a)
```

4. If everything went well, it should be possible to adopt the Makefile so that rewrite_system.h, rewrite_system.cpp can be included in the program. Run function test_rewrite( ), which uses the rewrite system

$$
\begin{array}{ll}
X+0 & \Rightarrow X \\
X+s(Y) & \Rightarrow s(X+Y) \\
X \times 0 & \Rightarrow 0 \\
X \times s(Y) & \Rightarrow(X \times Y)+X \\
E(X, X) & \Rightarrow t
\end{array}
$$

to test if $2 \cdot 2 \cdot 3$ equals $3 \cdot 2 \cdot 2$.

You can also make some other tests.

